



# On Multi-Gateway Association in Wireless Mesh Networks

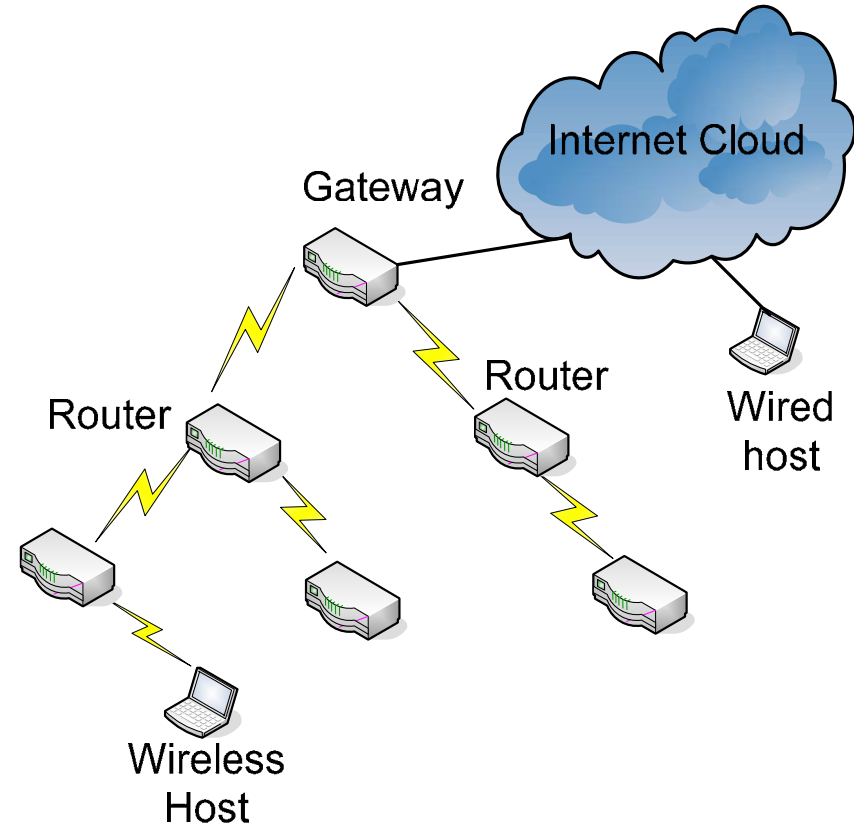
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# Wireless Mesh Networks

- Characteristics
  - Multi-hop with a static routing infrastructure
  - Several applications in military and commercial scenarios such as hospital, residential broadband networks
  - **Gateway Association** – establishing a connection to the wired internet( through a gateway)
- Challenges
  - capacity, security and fairness

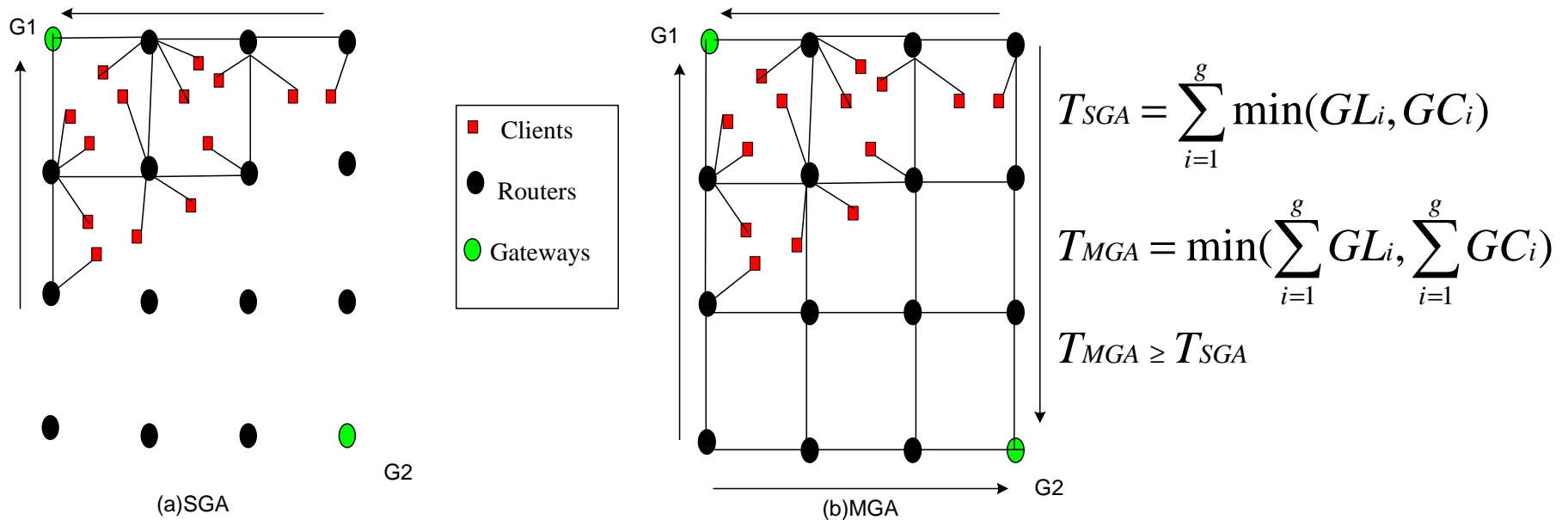


# Motivation

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- Default model
  - Single Gateway Association [Wimesh2005].
    - simple model
    - leads to several issues such as single-gateway bottleneck, fairness, security ;etc
  - What if we move away from this model?
- Proposed Model
  - Multi-Gateway Association (MGA)
    - every client can use more than one gateway simultaneously
- Focus of present work
  - highlight the benefits & challenges of MGA
  - develop a (centralized) solution suite to achieve anticipated benefits
  - evaluate solution performance using simulations

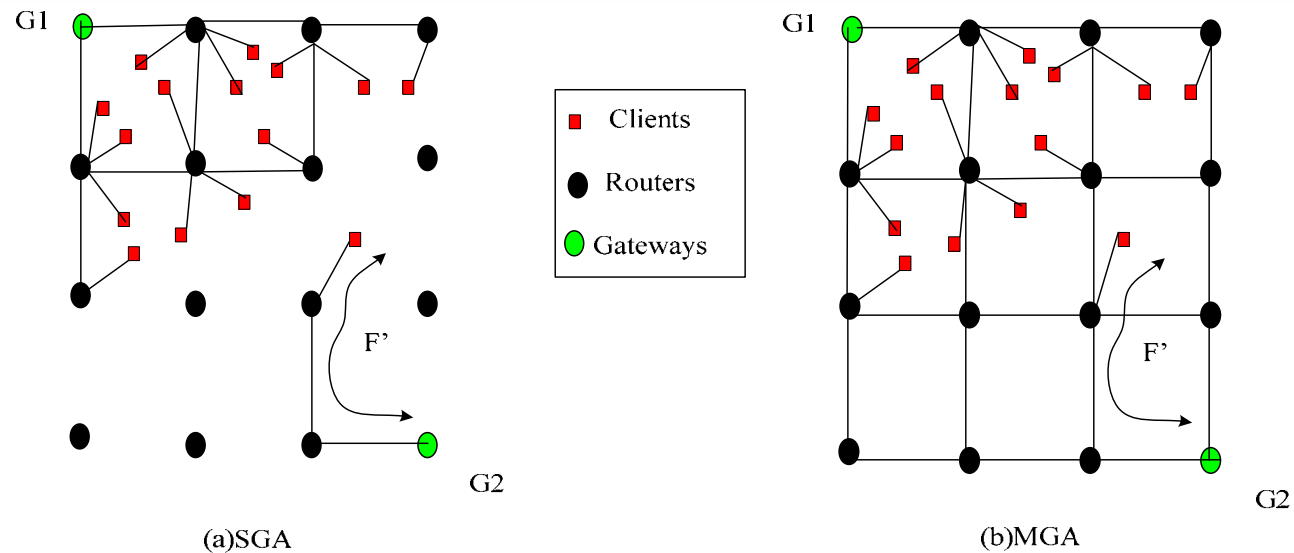
# Benefits of MGA-Capacity(1/4)



## ■ Capacity

- Better resource utilization
- Ideally when the network load is below throughput capacity

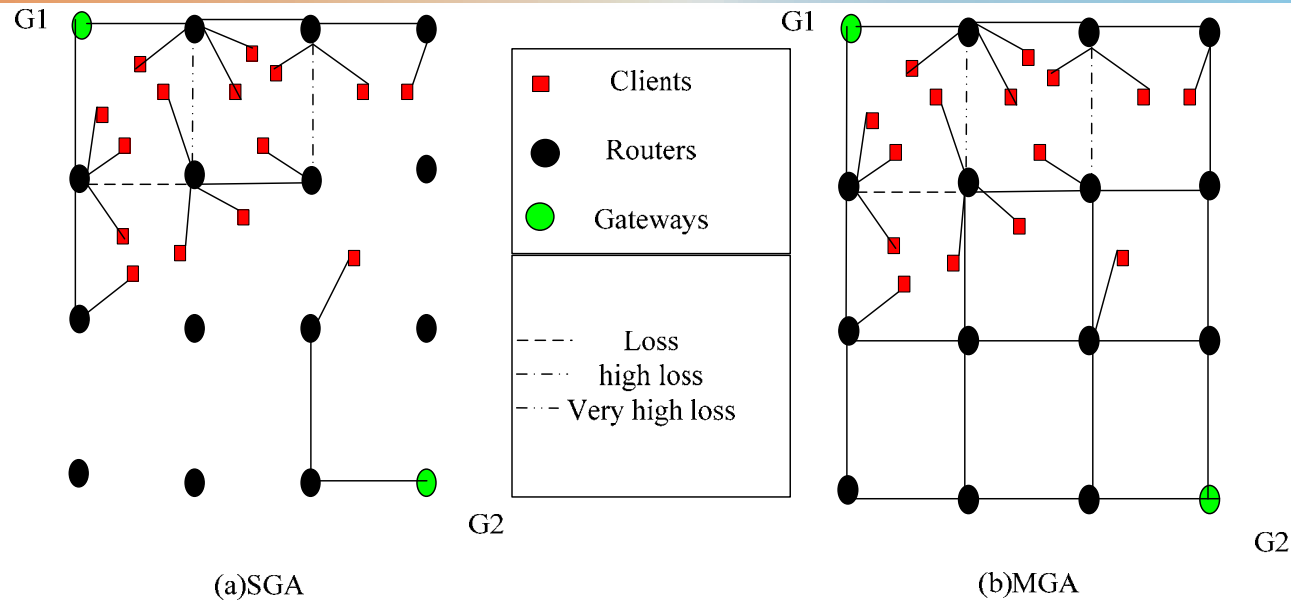
# Benefits of MGA-Fairness (2/4)



## ■ Fairness

- Problems due to
  - Uneven geographic distributions – nearness to gateway
  - Uneven traffic distribution – differently loaded gateways
- MGA improves fairness by more uniform distribution of loads and using aggregate resources than single gateway resources

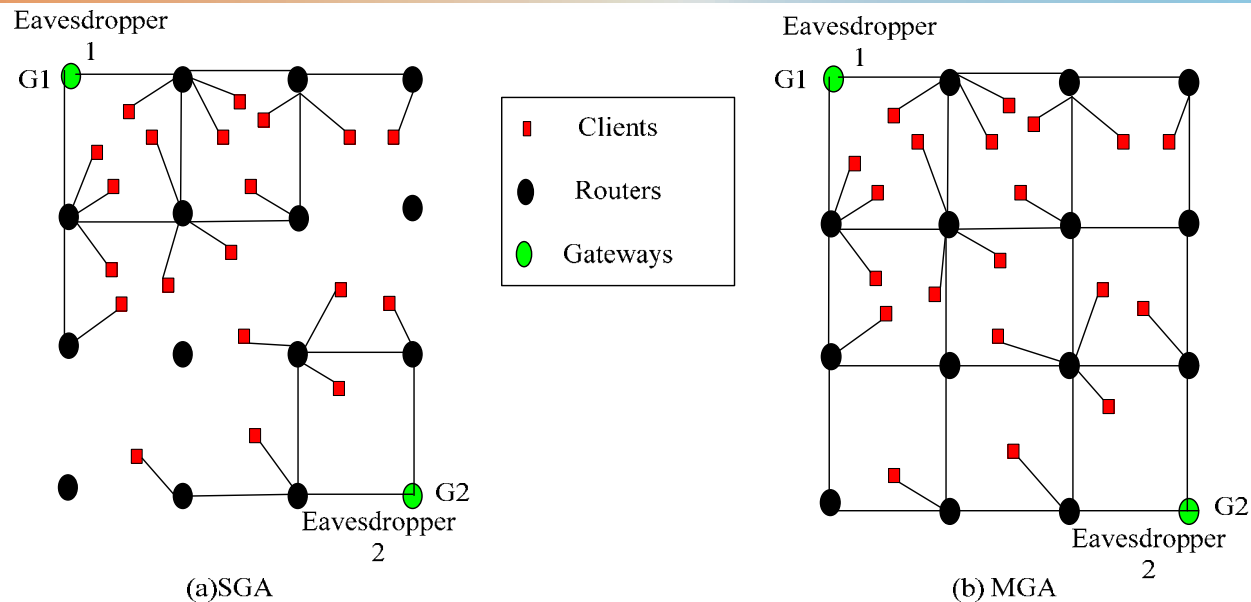
# Benefits of MGA-Diversity (3/4)



## ■ Diversity

- Types of loss:
  - Hard losses- failure of gateway
  - Soft losses- space dependent channel loss, buffer drops;etc
- Path, gateway diversity possible with MGA

# Benefits of MGA-Security (4/4)



## ■ Security

### ■ Problem:

- Gateway - a single aggregation point
- Interception around gateway exposes the wireless network

### ■ MGA

- Makes the eavesdropper's task more difficult
- Reduces the number of intercepted packets

# Challenges

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- **Architectural model**
  - Requirement of splitting for downstream and reassembly for upstream
- **Gateway characterization**
  - loss , delay and throughput characteristics of the paths from each client to each gateway
- **Gateway Association**
  - How many gateways to associate with?
  - Which gateways to associate with?
- **Scheduling**
  - which packet to send at what time instant so that effective aggregation and in-sequence delivery occur?



# Gateway Association Algorithm

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- Overview
  - Greedy centralized algorithm to determine associations and rates
  - Maximize aggregate throughput subject to max-min fairness
- Algorithm
  - Step 1: Path Computation
    - Compute shortest paths to each gateway from the client
  - Step 2: Bandwidth Computation
    - Compute available bandwidths on each path using the flow graph
  - Step 3: Gateway Decision
    - Identify the number of gateways and the exact gateways to associate with

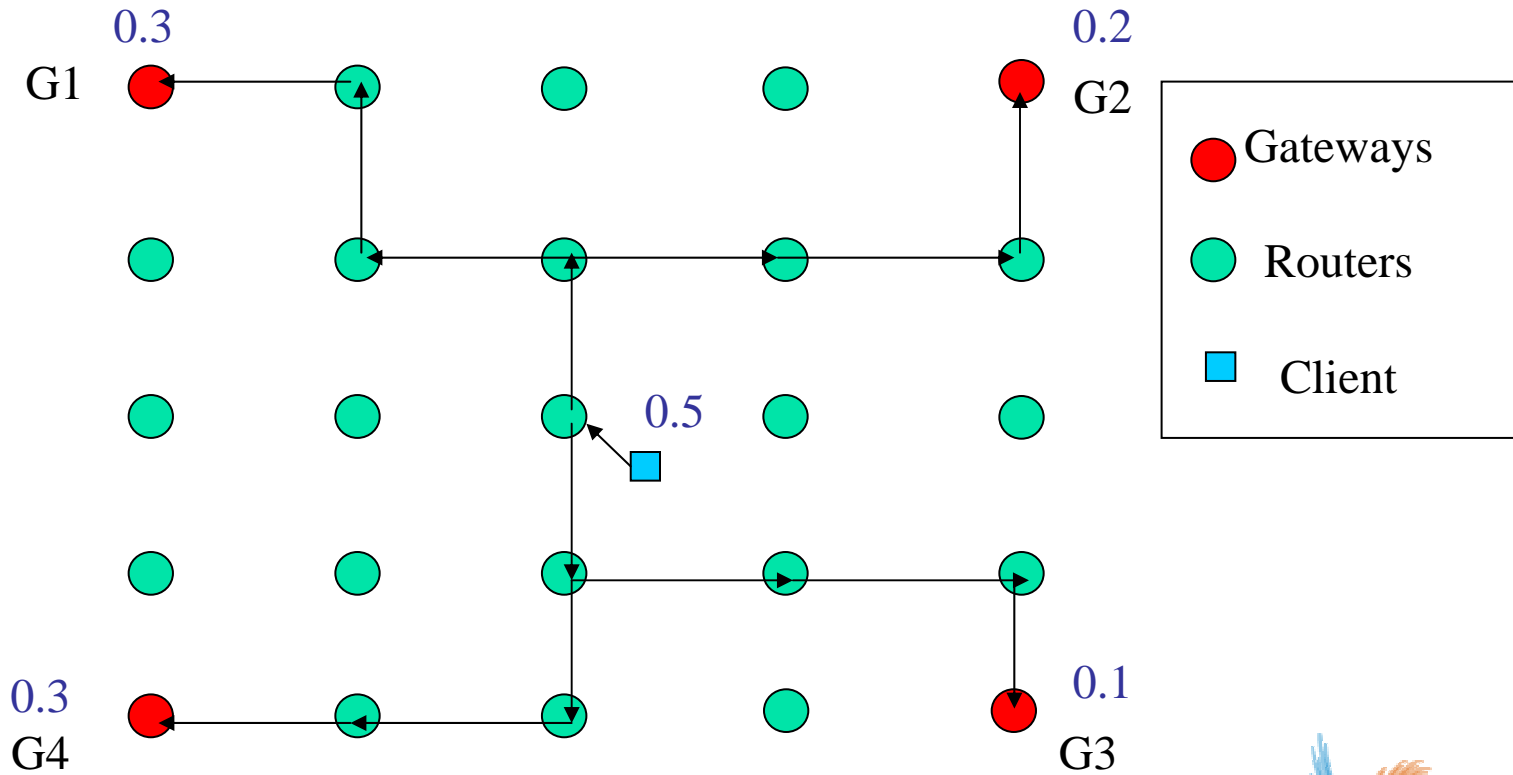
# Gateway Decision Details

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- Is load less than sum of available bandwidths?
  - Yes: allot bandwidth in a greedy manner (in descending order of available bandwidths)
  - No: For the excess load
    - Compute maximum degree bottlenecks on paths to gateways
    - Identify the set of flows which contends or intersects to the maximum
    - Allot unit bandwidth from maximum rate flow to the new flow one the maximum bottleneck path
    - Terminate if load is allotted or max-min allocation is reached

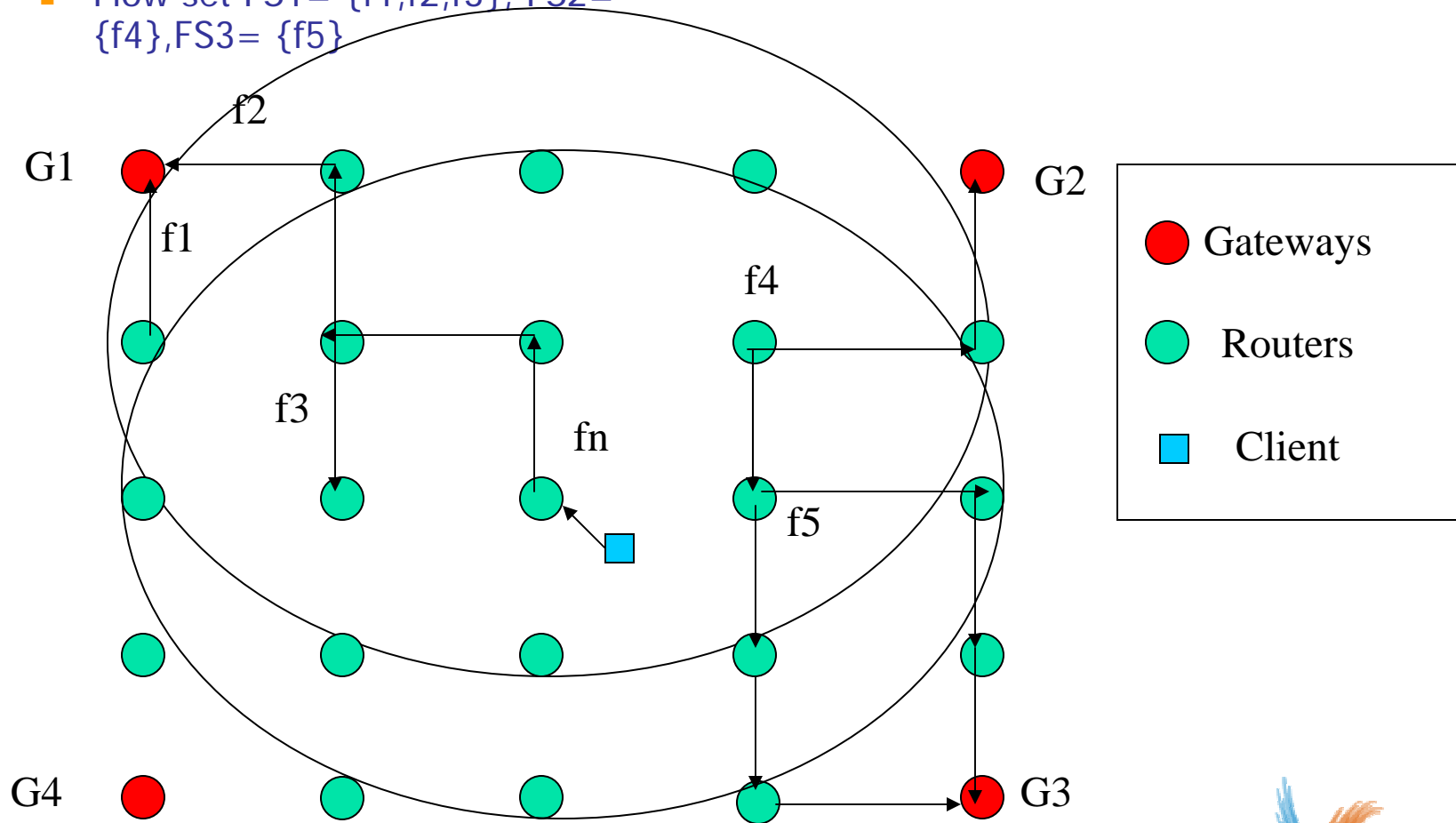
# Gateway Association Algorithm

- Illustration
  - Step 1 / Step 2
    - Client with load of 0.5 computes paths and identifies gateway statistics



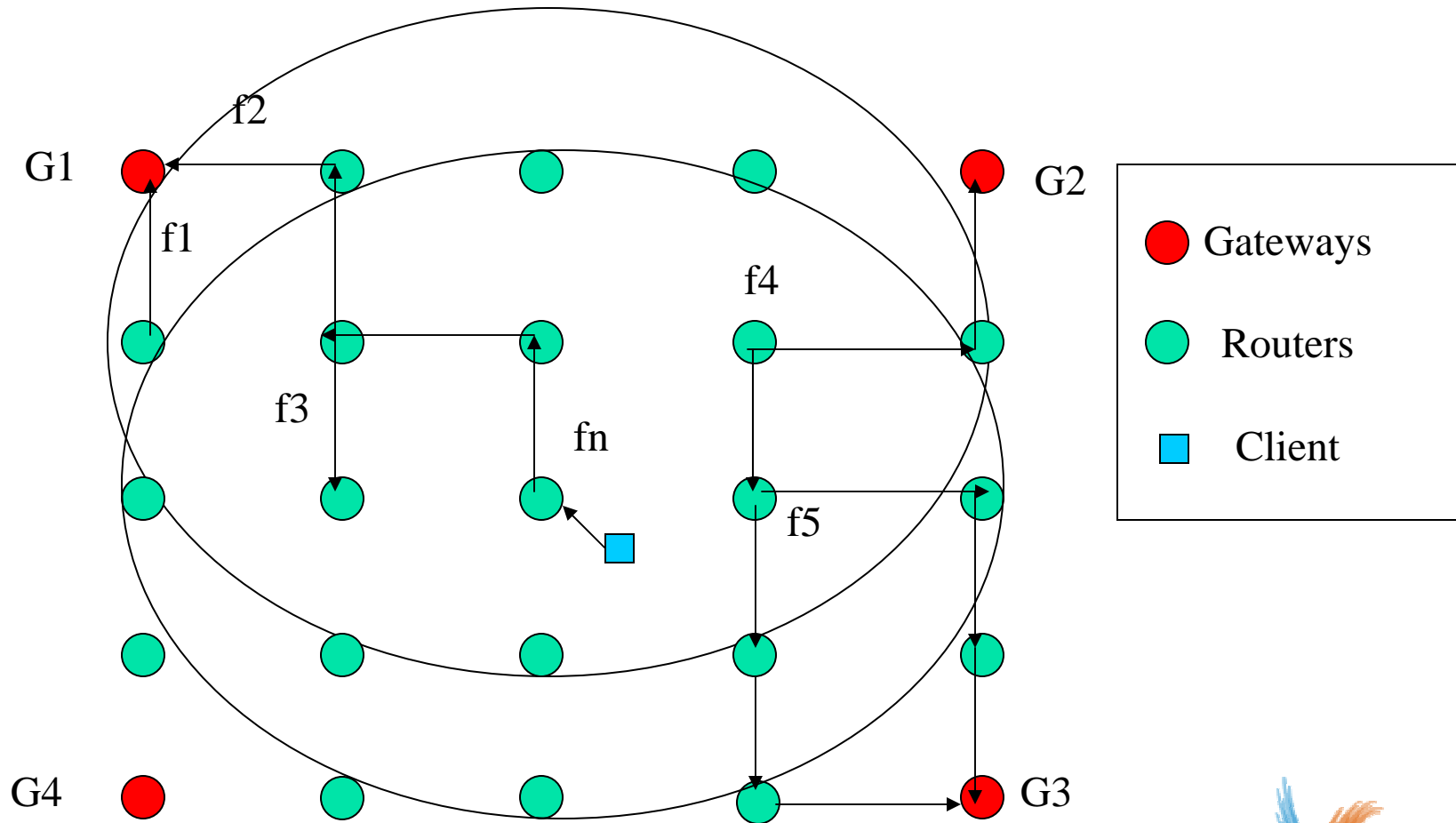
# GAA-Flow sets

- Illustration of step 3
  - Flow set  $FS1 = \{f1, f2, f3\}$ ,  $FS2 = \{f4\}$ ,  $FS3 = \{f5\}$



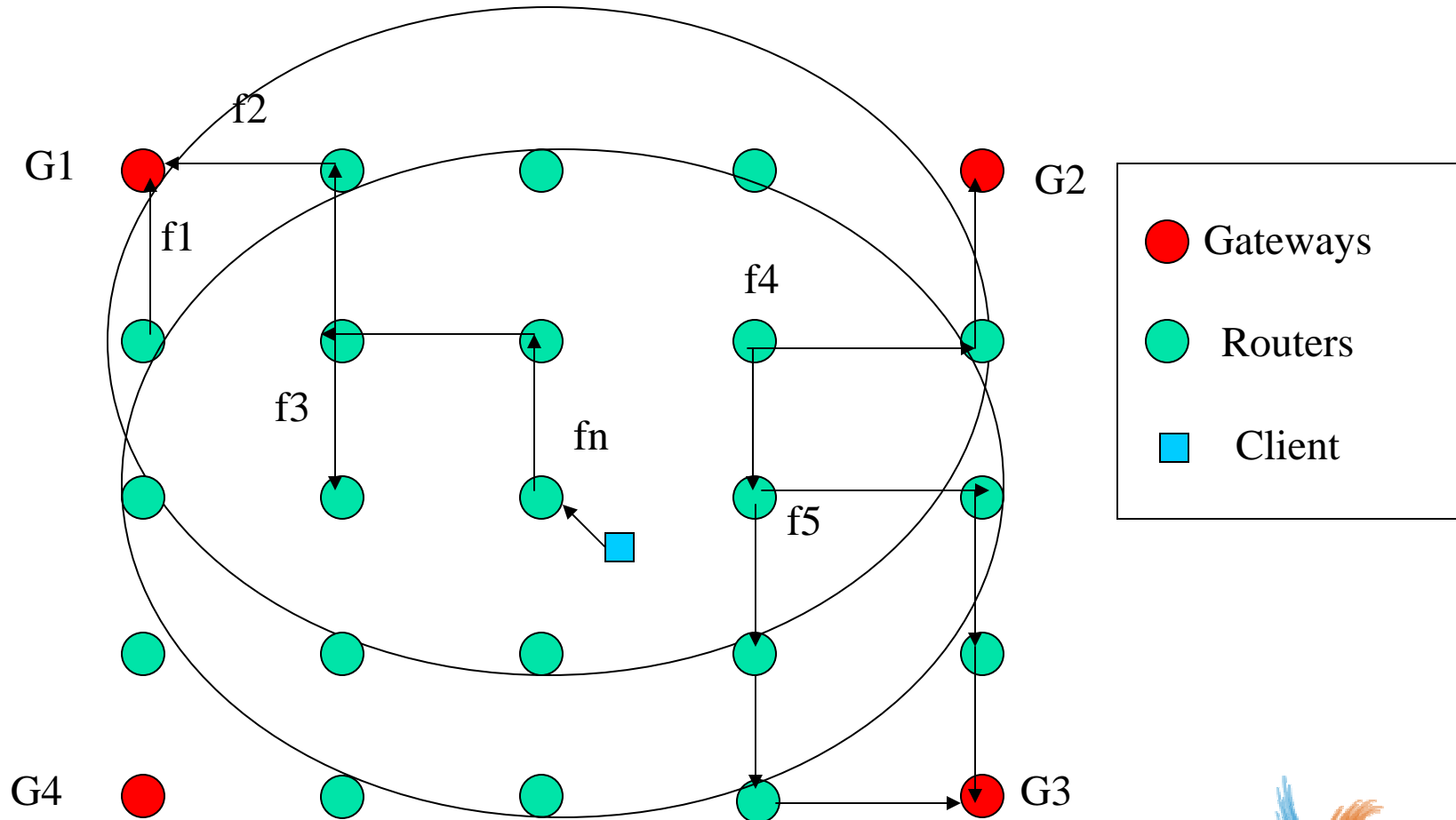
# Gateway Association Algorithm

OFS = {f1, f2, f3, f4, f5}



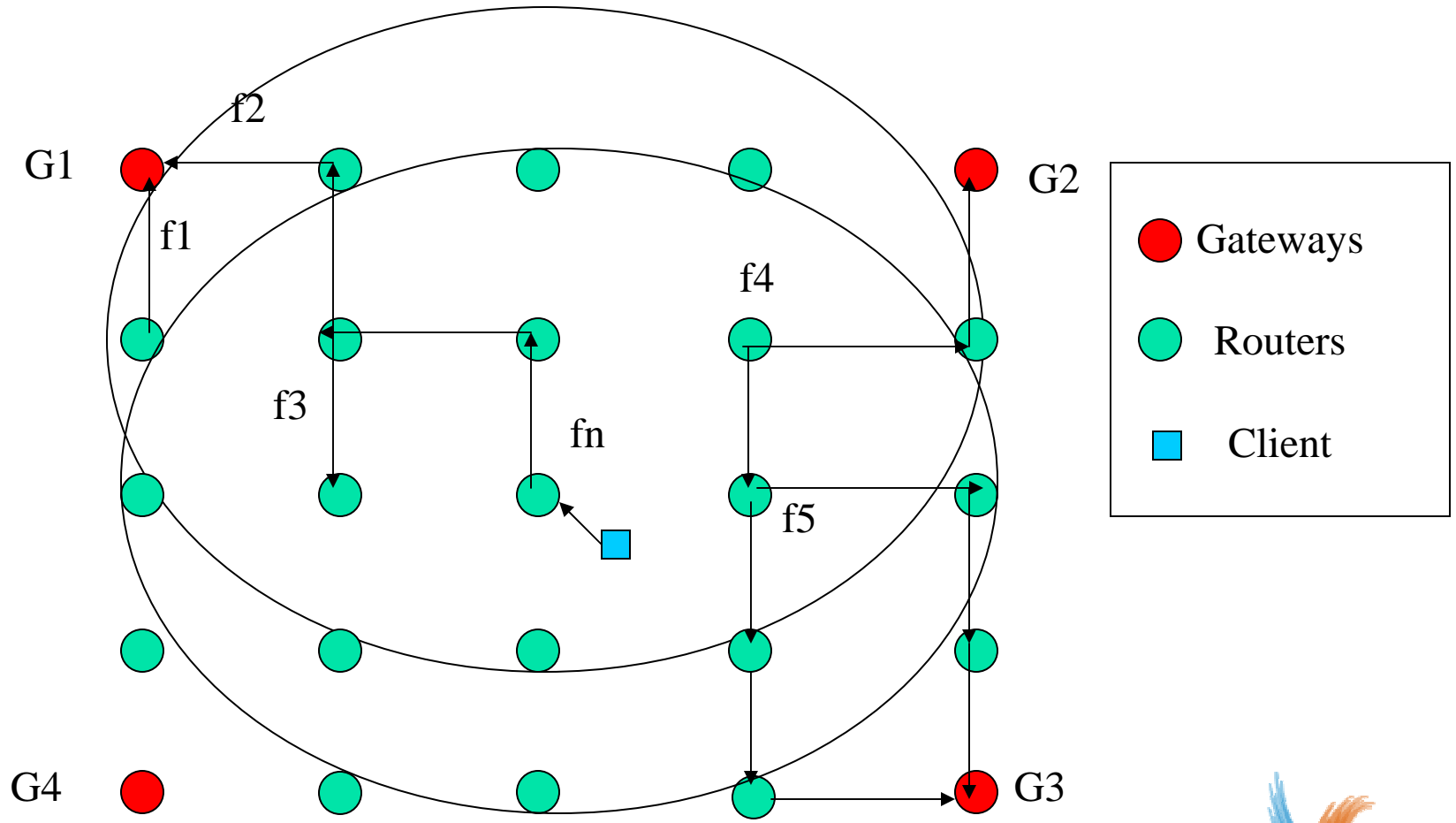
# Gateway Association Algorithm

- Let  $f_4$  be the maximum rate flow and  $\text{Rate}(f_4, 2) = 0.1$  ;  $\text{Rate}(f_4, 3) = 0.1$



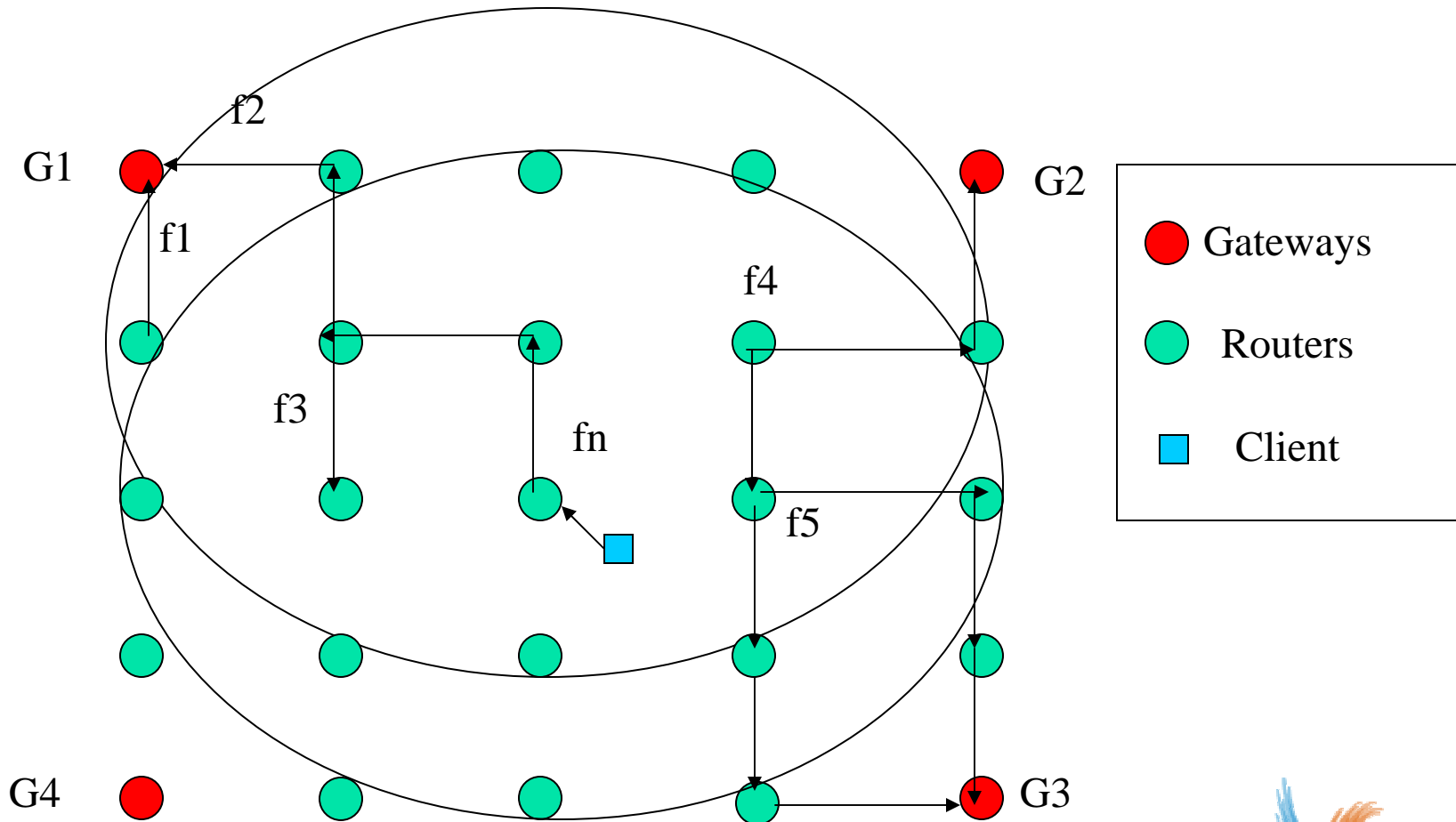
# Gateway Association Algorithm

- New rates:  $\text{Rate}(f_{4,2}) = 0.05$ ,  $\text{Rate}(f_{4,3}) = 0.1$ ;  $\text{Rate}(f_{n,2}) = 0.05$



# Gateway Association Algorithm

- Continues till max-min allocation or  $f_n$ 's requirement





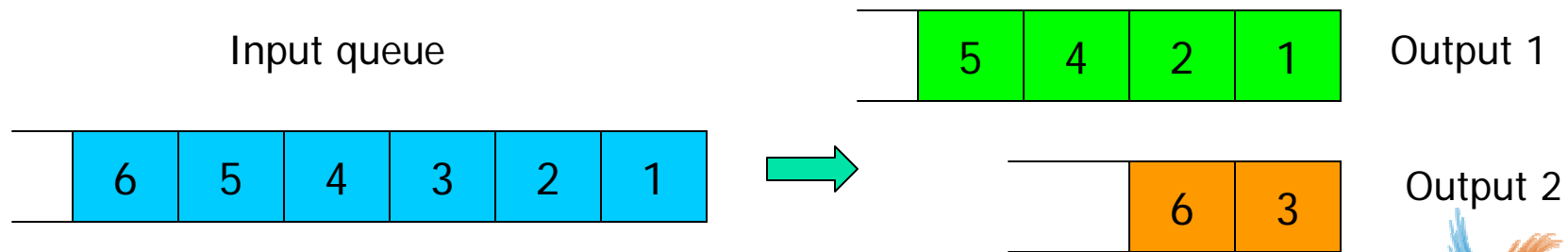
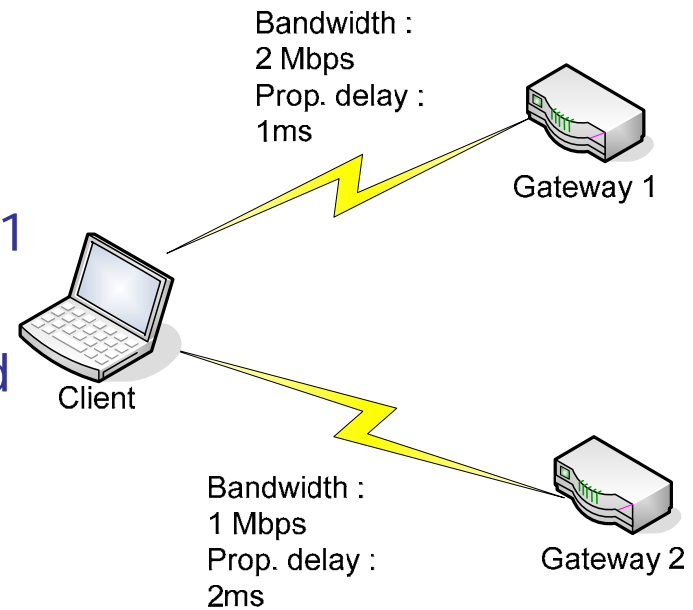
# Scheduling - Overview

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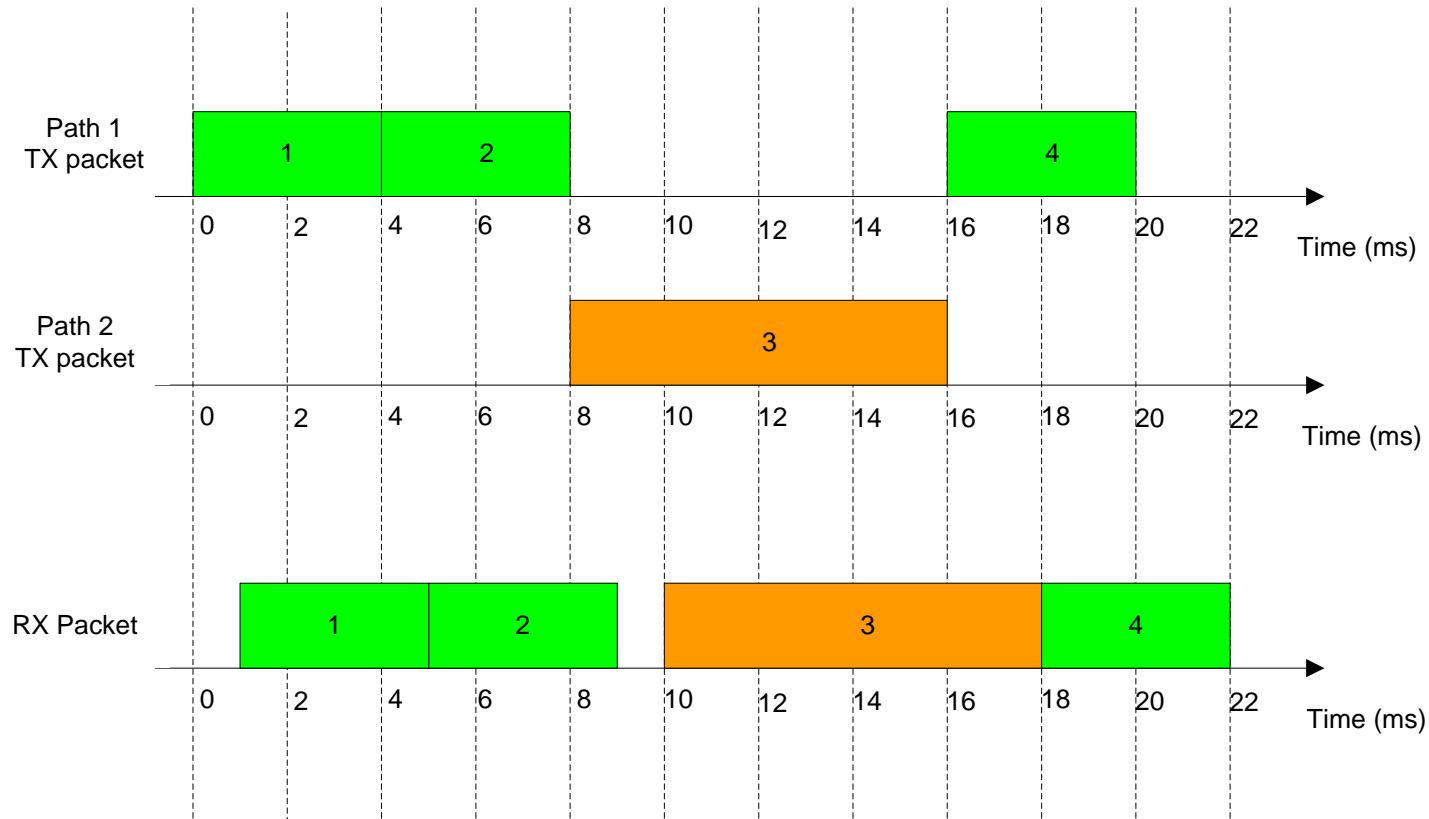
- Goal
  - enable effective aggregation and prevent out of order delivery
- Procedures
  - Step 1: Window calculation
    - Calculate the number of packets to each of the associated gateways
    - Uses integral values of normalized delay ratios.
  - Step 2: Ordering
    - Calculate reception times as transmission time + other delays
    - Tag each packet with a 2-tuple (seq.no, expected reception time)
      - Expected reception time for packets of the same destination = Propagation delay +  $k$ \*Transmission delay ( where  $k$  is its position in the destination queue)
    - Sort in ascending order of reception time

# Scheduling Illustration

- Window calculation
  - The total delays (Assuming a packet size of 1000 bytes) are
    - $D1 = 4 \text{ ms} + 1 \text{ ms} = 5 \text{ ms}$
    - $D2 = 8 \text{ ms} + 2 \text{ ms} = 10 \text{ ms}$
  - Normalized integral delays are 2 and 1
- Ordering
  - Generate 2-tuples (seq. no , expected reception time) for example ( 1,5)  
(2,9) (3,10)



# Scheduling Instants



- In-sequence delivery, if delays are correct

# Performance evaluation - Setup

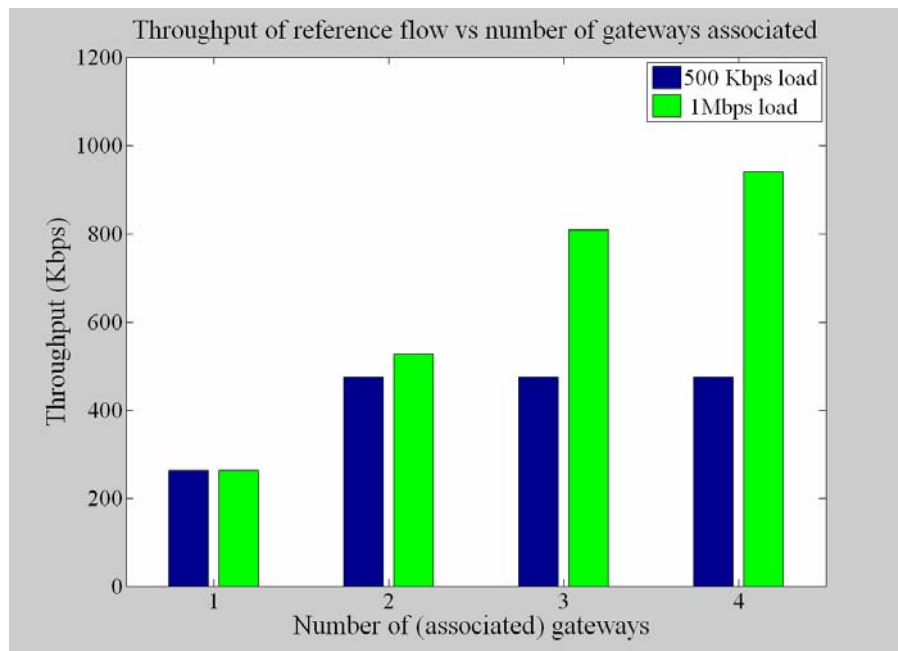
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- Ns2 simulator
- 1000 m \* 1000 m grid
- 4 gateways at (200,200) , (200,900) , (900,200) and (900,900)
- 21 Routers uniformly deployed along with 10 clients
- Each client has a traffic demand of 500 Kbps
- Application – CBR
- Transport layer- UDP
- Routing- Static Shortest path
- MAC - Ideal Flow Scheduling
- PHY – Wireless Phy at 915 MHz

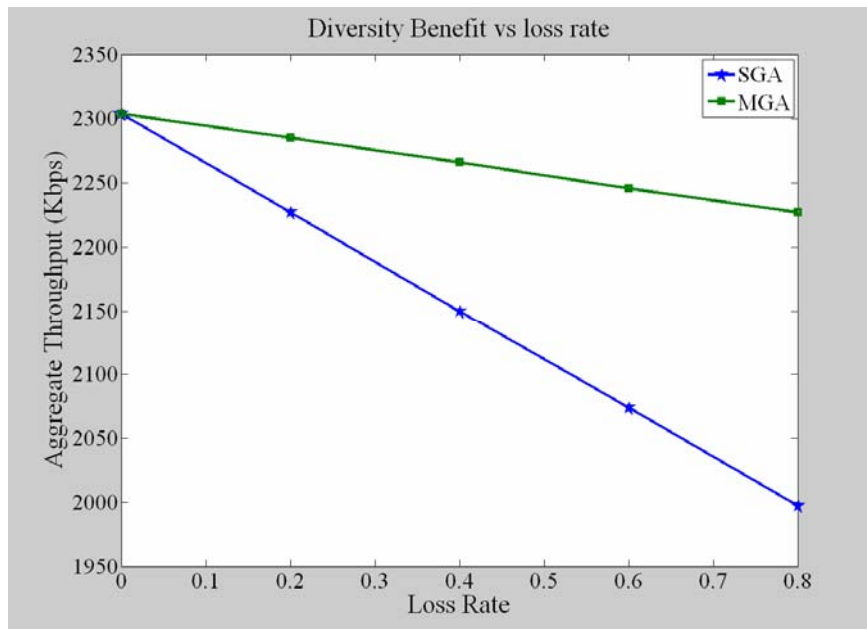
# Throughput Improvements

SGA	MGA	benefit
2.5Mbps	2.76 Mbps	10.4%

- Average case
  - improvement around 10% .
- Effect of load
  - Single reference client at the center and other clients outside the grid of gateways
  - Almost linear improvement in throughput gain with number of associations
  - Load also decides the gains possible

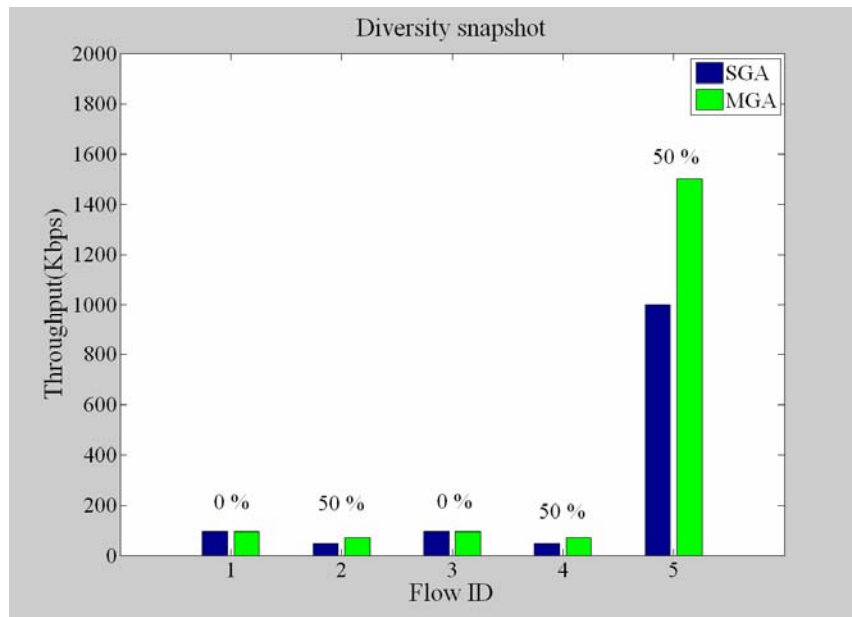


# Diversity Benefit (1/2)



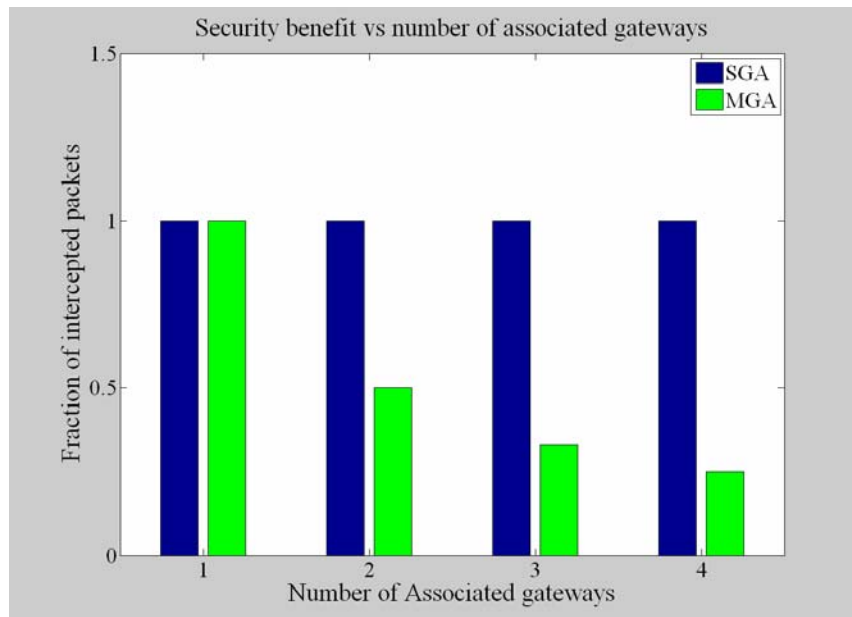
- A single loss module at gateway 0 ( i.e. at (200,200))
- Simulated path loss rate at gateway
- Averaged over 10 seeds for each loss rate
- MGA has lesser degradation and rate of degradation with losses

# Diversity Benefit (2/2)



- For a fixed 50% loss rate at one gateway
- Loads are 100,100,100,100, 2000 Kbps
- Flows achieve 50% benefit in throughput
- Only flows associated with lossy gateway show improvement

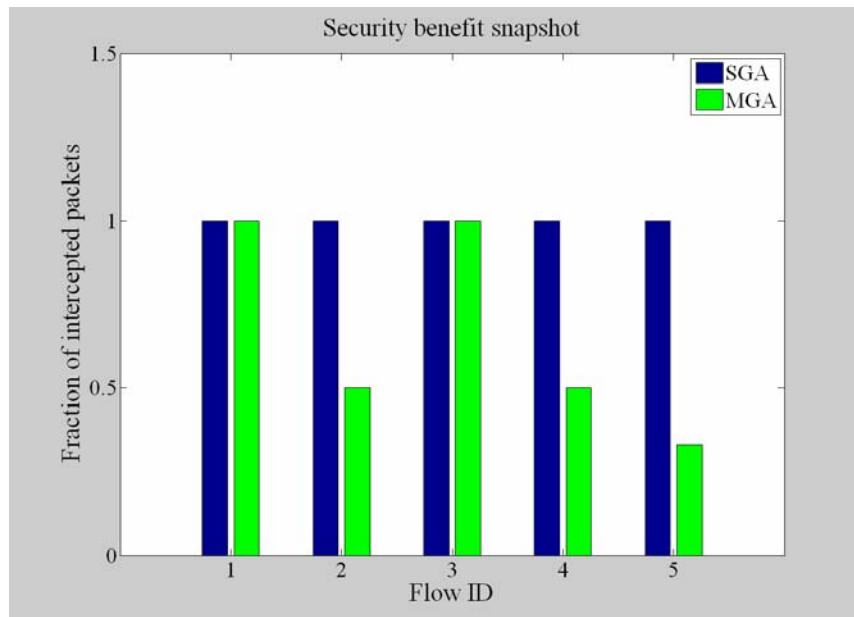
# Security Benefit (1/2)



- Security benefit increases with number of gateways
- Fraction of intercepted packets of each flow with increasing associations allowed
- Total interception by eavesdropper assumed
- Trend is same for all flows as the number of associations is same
- Near ideal benefit ratios



# Security Benefit (2/2)



- Practical case of, providing security without significant (<10%) capacity degradation
- The associations are 1,2,1,2,3 for flows 1 to 5 respectively
- Near ideal benefits obtained for flows that perform multi-association

# Related Work

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- Wireless Mesh networks
  - Dual association of clients with routers for optimizing broadcast load [Wimesh 2005]
  - Capacity of mesh networks [Mobicom 2005]
  - Security challenges in wireless networks [Mobicom 2001]
- Multiple Connections
  - pTCP : Transport solution suite for Managing multiple TCP connections [Mobicom 2002]
  - R<sup>2</sup>CP : Receiver Centric Transport Protocol allows aggregation of bandwidth [Mobicom 2003]

# Conclusions

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- Showed that the MGA model has many benefits
- Identified challenges in leveraging the benefits of multi-gateway association
- Designed centralized algorithms for gateway association and scheduling
- Evaluated the association algorithm through simulations
- Multi-gateway association has benefits, but leveraging them requires non-trivial algorithm development
- Further neither a single gateway nor all gateway association can bring the best benefits

# Future work

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- Evaluate the scheduling algorithm and the effect of TCP flows
- Identify all the drawbacks of MGA
- Evaluate the different factors that impact the achievable benefits
- Design gateway characterization algorithms to obtain network information and decide in a distributed manner
- Identify a more general framework in which associations are decided to support the desired levels of capacity, diversity and security simultaneously

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Thank You